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# An analytical approach to compute lower bounds of $\mu$ -values

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## ABSTRACT

The lower bounds of  $\mu$ -values for a family of square real and complex valued matrices are computed analytically. The proposed methodology consists of factorizing an admissible perturbation from a set of block diagonal matrices into a block diagonal matrix. The computation of the lower bounds of  $\mu$ -values is then carried out by computing the spectral radius and numerical radius for matrix under consideration. The lower bounds of  $\mu$ -value provide the conditions which guarantee the instability analysis of the linear feedback system.

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## 1. Introduction

The numerical approximation of eigenvalues corresponding to the family of matrices plays a vital role in science and engineering. For instance, the largest eigenvalue computed by the power method corresponding to Leslie matrix describes the long term growth rate of population. Besides, the vibration frequencies are described with the help of eigenvalues of matrices appearing in structural mechanics, and the eigenvalues are the measure of data variance that can be used for dimensional reduction in multivariate data analysis.

The eigenvalues are roots of the characteristic polynomial. The roots are computable by using iterative methods. The class of matrices governing from real systems possesses uncertainties and the computation of properties for such matrices in NP - hard (Nemirovskii, 1993; Braatz et al., 1994).

The  $\mu$ -value is a well known MATLAB Tool available in MATLAB Control Tool Box (Doyle, 1982; Safonov, 1982; Safonov and Doyle, 1984) and has been used to discuss the stability, instability, performance, and robustness of feedback system in linear control. The numerical methods (Packard et al., 1988; Fan and Tits, 1986; Packard and Doyle, 1993) are used to approximate bounds of  $\mu$ -value but unfortunately the computation of its exact value is NP - hard (Braatz et al., 1994).

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The class of positive matrices like level symmetric matrices and Hermitian positive definite matrices are widely used in mathematics and in various applications of engineering. For instance, computer vision (Nemirovskii, 1993), the mechine learning (Kishida and Braatz, 2014) and in the area of convex optimization (Braatz et al., 1994). The Linear Matrix Inequality (LMI) technique based on the positive definiteness nature of matrix is widely used to study the stability analysis of feedback systems in linear control. In control, however, various control systems are designed on the top of state-space models, which are symmetric in nature; this includes power networks and an electrical network (Doyle, 1982).

The algorithm provides tighter lower bounds for  $\mu$  –values when real uncertainties are under consideration (Dailey, 1990). The proposed algorithm is based on simple operation of matrix algebra, and it iterates with respect to only one variable, which returns not only the size of the worst case parameter but also its actual values. In Karamancıoğlu and Kasimbeyli (2011), a non-linear programming technique is introduced to approximate the tighter lower bounds of real  $\mu$  –values. The real structured singular value problem (RSSV) is formulated as a non-linear programming problem, which is then solved by making use of the F-modified sub-gradient (F-MSG) technique to compute lower bounds of structured singular value. The F-MSG algorithm solves a large class of non-convex optimization problems without making use of differentiability.

In Kim et al. (2009), a geometrical approach is introduced to approximate the lower bounds of  $\mu$  –values for pure real repeated perturbations. The problem formulation appears in the sense that the resulting parametric search space does not depend

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on the fact that how many times the parameter is repeated in a structured perturbation matrix.

In Fabrizi et al. (2014), a detailed comparison of the developed numerical method is presented. Each listed numerical method in Fabrizi et al. (2014) approximates and gives improved results for lower bounds of  $\mu$  –values.

A Gain-Based lower bound algorithm is presented in Seiler et al. (2010); to compute lower bounds of  $\mu$ -values. The Gain-Based lower bound algorithm takes both real and mixed perturbation along with the given matrix whose  $\mu$ -values computations are under consideration. The key idea of this algorithm is to make use of worst-case gain problems in order to approximate the real perturbation, and it uses a standard power algorithm to compute the complex perturbation.

In Rehman et al. (2019), the lower bounds of  $\mu$  –values are approximated by making use of low rank ODEs based technique. The proposed methodology is based on two level algorithm, the inner-outer algorithm. The proposed iterative method approximated tighter lower bounds of  $\mu$  –values when compared with the well-known MATLAB function MUSSV in most cases.

This manuscript is organized as Section 2 consists of mathematical preliminaries, which includes a basic definition of the spectrum, pseudo-spectra, and  $\mu$  –values. In Section 3 we present our methodology, which consists of the decomposition of uncertainty from a set of block diagonal matrices into a block diagonal matrices that have identities matrices along with principle diagonal. Section 4 summarizes our findings.

#### 2. Mathematical preliminaries

The real and complex *n* dimensional matrices are denoted by  $K^{n,n}$  with K = C (*or* R). The real and complex scalars are denoted by K = C (*or* R). The complex vectors are denoted by  $C^{n}$ .

For a complex matrix M,  $M^*$  denotes the complex conjugate transpose. The  $n \times n$  identity matrix is denoted by  $I_n$  with n representing its dimensions. The spectrum of a matrix is represented with notation  $\Lambda(\cdot)$  while  $\Lambda_{\epsilon}(\cdot)$  represents the  $\epsilon$ -pseudospectrum of a matrix (·). The notation  $\|\cdot\|$  represents the norm of a matrix or vector. The notations **B** and **B**<sup>\*</sup> denote the set of block diagonal matrices having mixed real and complex uncertainties and pure complex uncertainties, respectively.  $\rho(\cdot)$  denotes the spectral radius of a matrix (·).

**Definition 2.1:** The set of eigenvalues of an ndimensional complex matrix  $M \in \mathbf{K}^{n,n}$  with  $\mathbf{K} = \mathbf{C}$  is defined as:

$$\Lambda(M) = \{\lambda \in \mathbf{C} : det(M - \lambda I) = 0\}.$$
(1)

**Definition 2.2:** The  $\epsilon$ -pseudospectrum of a complex valued matrix  $M \in C^{n,n}$  with a small parameter  $\epsilon > 0$  is defined as.

$$\Lambda_{\epsilon}(M) := \{\lambda \in \mathbf{C} : \left| \frac{1}{(M - \lambda I)} \right| \ge \frac{1}{\epsilon} \}.$$
<sup>(2)</sup>

**Definition 2.3:** For  $\epsilon > 0$ , a scalar  $\lambda \in C(\text{or } \mathbf{R})$  belongs to the  $\epsilon$ -pseudo-spectrum of  $M \in \mathbf{C}^{n,n}$  and satisfies the following properties.

1.  $\lambda \in \Lambda(M + E)$  where  $\epsilon ≥ || E ||$  with *E* representing some perturbation to *M*.

2.  $\exists v \in \mathbf{C}^{\mathbf{n},\mathbf{1}}$  such that ||v|| = 1 and  $\epsilon \ge ||Mv - \lambda v||$ . 3.  $\lambda \in \rho(M)$  and  $\epsilon^{-1} \le ||\frac{1}{(M-\lambda I)}||$ .

**Definition 2.4:** The set of block diagonal matrices with mixed real and complex uncertainties is defined by:

$$\mathbf{B} = \{ diag(\delta_i I_i, \Delta_j) : \delta_i \in \mathbf{K}, \Delta_j \in \mathbf{K}^{m_j, m_j} \}.$$
(3)

**Definition 2.5:** The set of block diagonal matrices with pure complex uncertainties is defined as:

$$\mathbf{B}^* = \{ diag(\delta_i I_i, \Delta_j) : \delta_i \in \mathbf{C}, \Delta_j \in \mathbf{C}^{m_j, m_j} \}.$$
(4)

**Definition 2.6:** For  $M \in C^{n,n}$  and **B**, the  $\mu$  –value is defined as:

0 if $det(I - M\Delta) \neq 0$ , $\forall \Delta \in \mathbf{B}$		
$\mu_{\mathbf{B}}(M) = \{(\min\{\ \Delta\ _2: det(I - M\Delta) = 0, \forall \Delta \in \mathbf{B}\})^{-1},$	elsewhere, (5)	

where  $\|\Delta\|_2$  denotes the largest singular value of an admissible perturbation  $\Delta \in \mathbf{B}$  and  $det(\cdot)$  denotes the determinant of matrix (·).

**Definition 2.7:** For  $M \in \mathbf{C}^{n,n}$  and  $\mathbf{B}^*$ , the  $\mu$ -value is defined as:

$$0 \quad \text{if} \quad \rho(M\Delta) \neq 1, \quad \forall \Delta \in \mathbf{B}^*$$
  
$$\mu_{\mathbf{B}^*}(M) = \{ (\min_{\Delta} \{ \| \Delta \|_2 : \rho(M\Delta) = 1, \quad \forall \Delta \in \mathbf{B}^* \} )^{-1}, \text{ elsewhere.}$$
(6)

## 3. Computing structured singular values

In this section, we present an analytical approach for the computation of lower bounds of  $\mu$  –values. The proposed methodology is based on the idea of factorizing the set of a block of diagonal uncertainties into diagonal matrices having identity matrices along principal diagonal. Furthermore, the proposed methodology involves the computation of spectral and numerical radii of a given matrix, which in turn computes the lower bounds of  $\mu$  –values.

**Definition 3.1:** The modified block diagonal matrix is denoted by B and is defined as:

$$B = diag(A_1, A_2, \dots, A_n) diag(A_1^{-1}, A_2^{-1}, \dots, A_n^{-1}) = diag(I_1, I_2, \dots, I_n),$$

with each  $I_i \forall i = 1: n$  possess the same dimension. **Definition 3.2:** The spectrum of a matrix M is denoted by  $\sigma(M)$  and is defined as:

$$\sigma(M) = \{\lambda_i : \lambda_i \text{ is a neigenvalue}\}.$$
(7)

**Definition 3.3:** A matrix M is called a Stochastic matrix iff bM = b with b being a row vector of once. The vector  $b^t$  is an eigenvector of  $M^t$  corresponding to an eigenvalue  $\lambda \in \sigma(M)$  with  $|\lambda| = 1$ .

Let  $\Gamma^n$  be a space of *n* dimensional complex matrices and let  $H^n$  be a space of *n* dimensional Hermitian matrices. Let  $C^n$  be a unit vector.

**Definition 3.4:** The numerical range W(A) for  $A \in \Gamma^n$  is defined as:

$$W(A) = \{u^* A u : u \in \mathbf{C}^n, \| u \|_2 = 1\}.$$
(8)

**Definition 3.5:** The numerical radius r(A) of  $A \in \Gamma^n$  is defined as:

$$r(A) = \max\{Re(v^*u^*Au): v \in \mathbf{C}, ||v|| = 1, ||u||_2 = 1 \text{ for } u \in \mathbf{C}^n\},$$
(9)

where  $v \in \mathbf{C}$  with || v || = 1 is a scalar on the unite circle  $\hat{\mathbf{C}}$ .

For  $B = diag(\hat{\delta}_1 \hat{A}_1, \hat{\delta}_2 \hat{A}_2, \dots, \hat{\delta}_n \hat{A}_n)$  we assume that each  $\hat{\delta}_i \hat{A}_i, \forall i = 1: n, \hat{\delta}_n \in \mathbf{C}(\mathbf{R}), \hat{A}_i \in \mathbf{C}^{\mathbf{n},\mathbf{n}}(\mathbf{R}^{\mathbf{n},\mathbf{n}})$ are invertible matrices. From here and onwards, we shall write  $\hat{\delta}_i \hat{A}_i = A_i \forall i = 1: n$ . Theorem 3.6 computes a block diagonal matrix having identity matrices along its principal diagonal. This idea helps in order to compute the spectral and numerical radii for a given matrix rather than computing these mathematical quantities for the product of a given matrix with admissible perturbation.

**Theorem 3.6:** Let A and B be n-dimensional square matrices which are identically partitioned into block diagonal matrices:

$$\begin{aligned} A &= diag(\delta_1 A_1, \delta_2 A_2, \cdots, \delta_n A_n), \quad B = \\ diag(\delta_1 B_1, \delta_2 B_2, \cdots, \delta_n B_n), \quad \delta_i \in \mathbf{C}(\mathbf{R}) \quad \forall \ i = 1:n. \end{aligned}$$
 (10)

The product

$$AB = diag(\delta_1 A_1, \delta_2 A_2, \cdots, \delta_n A_n) diag(\delta_1 B_1, \delta_2 B_2, \cdots, \delta_n B_n)$$

is a block diagonal matrix  $C = diag(\delta_1 C_1, \delta_2 C_2, \dots, \delta_n C_n)$ , which is identically partitioned to block diagonal matrices. Thus in *C* each  $C_k = A_k B_k$   $\forall k = 1: n$ .

**Proof:** As block diagonal matrices  $diag(\delta_i A_i) \forall i = 1:n$  and  $diag(\delta_i B_i) \forall i = 1:n$  are identically partitioned, which in turn implies that each of these block diagonal matrices possesses k number of columns and k number of rows. Partitioning of block diagonal matrices can be presented while making use of some non-zero vector  $\vec{u}$  having the length k + 1. In such a situation  $A_k, B_k \forall k = 1:n$  possesses  $(u_{k+1} - u_k)$  number of non-zero columns and rows. Since the  $diag(\delta_i A_i)diag(\delta_i B_i) =$ 

 $diag(\delta_i C_i) \forall i = 1:n$  which means that  $diag(\delta_i C_i) \forall i = 1:n$  can be described in terms of block diagonal matrix, which possesses k number of non-zero columns and rows. The matrix C is partitioned as  $C = [C_1, C_2, \dots, C_k]$  with  $C_1, C_2, \dots, C_k$  being k columns of C. The submatrices  $C_{k\beta}$  in C, takes the form,

$$C_{k\beta} = \sum_{r=1}^{k} A_{kr} B_{r\beta}.$$
 (11)

In turn, this implies that,

$$C_{k\beta} = \sum_{r=1}^{k-1} A_{kr} B_{r\beta} + A_{kk} B_{k\beta} + \sum_{r=k+1}^{k} A_{kr} B_{r\beta} \quad \text{with} \quad r \neq k.$$
(12)

As  $A_{kr} = 0$  is a null-matrix. This is because the matrix *A* is a block diagonal matrix. Thus,

$$C_{k\beta} = A_{kk} B_{k\beta} = 0. \tag{13}$$

On the other hand  $B_{k\beta} = 0$  as  $k \neq \beta$ . This shows that,

$$C_{kk} = A_{kk} B_{k\beta} \quad \text{or} \ C_k = A_k B_k, \tag{14}$$

it indicates that,

$$diag(\delta_i C_i) = dig(\delta_i A_i) diag(\delta_i B_i), \ \forall \ i = 1:n.$$
(15)

We give the following Theorem 3.7 in order to compute the spectral radius of a given matrix *M*. Furthermore, Theorem 3.7 shows that  $\lambda \in \sigma(M)$  attains the maximum value to be exactly equal to 1.

**Theorem 3.7:** Let M be a positive and diagonalizable matrix. Let  $\lambda \in \sigma(M)$  be an eigenvalue of M belonging to it's spectrum, that is,  $\sigma(M)$  of M. Consider  $|\lambda| = 1$  and Re( $\lambda$ ) > 0 then spectral radius,  $\rho(M) = 1$ .

**Proof:** We proceed with the proof by contradiction. Consider that  $Re(\lambda) > 0$ , that is, non-positive. Take  $\epsilon > 0$ , a small parameter such that  $diag(M) > \epsilon$ . Then, the eigenvalue,

$$|\lambda - \epsilon| > 1. \tag{16}$$

Take  $\delta \in (0,1)$  such that the eigenvalue  $\delta |\lambda - \epsilon|$  of matrix  $M_1 = \delta(M - \epsilon I)$  with *I* being an identity matrix such that,

$$\delta|\lambda - \epsilon| > 1 \tag{17}$$

and,

$$\rho(M_1) > 1. \tag{18}$$

Let  $M_2 = \delta M$ , then  $\rho(M_2) < 1$ . By making use of the fact that,

for 
$$\rho(A) < 1$$
,  $\lim_{n \to \infty} ||A^n|| = 0$ , (19)

and,

for 
$$\rho(A) > 1$$
,  $\lim_{n \to \infty} ||A^n|| = \infty$ , (20)

we have,

 $\lim_{n \to \infty} \| M_1^n \| = \infty, \text{ for } \rho(M_1) > 1$ (21)

and,

$$\lim_{n \to \infty} \| M_1^n \| = 0, \text{ for } \rho(M_2) < 1.$$
(22)

This indicates that  $M_1 > 0$  and  $(M_1^n)_{ij} \le (M_2^n)_{ij} \forall i, j, n$ . This is clearly a contradiction, thus  $\rho(M) = 1$ . Theorem 3.8 indicates that the algebraic multiplicity of a non-negative real valued Stochastic matrix is 1, that is,  $alg_m(max|\lambda_i|) = 1$ .

**Theorem 3.8:** Let M be a non-negative real valued Stochastic matrix while for some power k, M > 0, that is, the matrix M is positive. Then  $\exists$  a unique eigenvector w > 0 with,

 $Aw = w, \tag{23}$ 

such that  $\rho(A) = 1$  and for some vector  $b = (1,1,\dots,1)$ , bw = 1.

**Proof:** Let  $\lambda_i \in \Lambda(A)$  with  $\Lambda(A) = \{\lambda_i: \text{Eigenvalues of } A\}$  and we aim to show that  $alg_m(max|\lambda_i|) = 1$ . For this, we consider that  $|\lambda_{max}(A)| \in \Lambda(A)$  is zero. But for  $A \ge 0$  one can have  $A^k = 0$  for some k. This leads to a contradiction and implies that  $|\lambda_{max}(A)| \in \Lambda_o(A)$  with  $\Lambda_o(A) = \{|\lambda_i: \lambda > 0 \quad \forall i|\}$ . As a consequence  $max_i|\lambda(i)| = 1$ .

Suppose that  $alg_m(A) > 1$ . The Jordan canonical form of matrix  $A \ge 0$  implies that  $\exists \quad \vec{v} \ne 0$  such that,

$$A\vec{v} = \vec{v} \tag{24}$$

or,

$$A\vec{v} = \vec{v} + \vec{z}$$
 for some  $\vec{z}$  corresponding to  $max_i|\lambda(i)| = 1.$ 
(25)

Theorem 3.9 shows that an Eigen spectrum of a Hermitian matrix contains the maximum eigenvalue to be exactly equal to 1, that is,  $max_i|\lambda_i(A)| = 1$ .

**Theorem 3.9:** Let  $T: \Gamma^n \to H^n$  with  $T(A) = \frac{1}{2}(A + A^*)$ and  $A \in \Gamma^n$ . Then  $\lambda_{max}(\frac{1}{2}v^*(A + A^*)) = 1$  and is equivalent to,

$$det(\frac{1}{2}v^*(A+A^*) - I) = 0, \quad \forall v \in \hat{\mathbf{C}},$$
(26)
where *I* is an identity matrix.

**Proof:** We proceed with the proof with contradiction. Suppose that,

$$det(\frac{1}{2}v^{*}(A+A^{*})-I) \neq \lambda_{max}(\frac{1}{2}v^{*}(A_{i}+A_{i}^{*})) = 1.$$
(27)

This means that  $\exists$  a sequence of matrices  $A_i \in \Gamma^n$ and these matrices converge to say  $\tilde{A} \in \Gamma^n$ , also  $\exists v_i \in \hat{C} \forall i$  such that,

$$\tilde{\lambda} \in \Lambda(\frac{1}{2}v_i(A_i + A_i^*)), \tag{28}$$

with  $\tilde{\lambda} = 1$  and  $\Lambda(0)$ , the spectrum of  $(\frac{1}{2}v^*(A + A^*))$ . But,

$$\tilde{\lambda} \neq \lambda_{max}(\frac{1}{2}\nu^*(A+A^*)).$$
<sup>(29)</sup>

The eigenvalue  $\lambda_2$  is the second largest eigenvalue that satisfies the fact,

$$\lambda_2(\frac{1}{2}v^*(A+A^*)) \ge \tilde{\lambda}(\frac{1}{2}v^*(A_i+A_i^*)) \quad \forall \ i.$$
(30)

But, for the subsequence we take  $v_i \rightarrow v \in \hat{\mathbf{C}}$  and,

$$\lambda_2(\frac{1}{2}v^*(\tilde{A}+\tilde{A}^*)) \ge \tilde{\lambda}(\frac{1}{2}v^*(A_i+A_i^*)),$$
(31)

which is a contradiction to our assumption which ensures that  $(\frac{1}{2}v^*(A_i + A_i^*))$  possesses an eigenvalue  $\tilde{\lambda} = 1$  as the second largest eigenvalue.

The eigenvector  $\vec{z} = \vec{v} + \alpha \vec{v} > 0$  for some  $\alpha \in R^+$ . Thus,

$$[A\vec{z}]_i = [\vec{z} + \vec{v} + \alpha \vec{v}]_i > \vec{v}_i \quad \forall i.$$
(32)

This is clearly a contradiction and hence  $alg_m(A) = 1$ .

Next, we show that for Stochastic matrix  $A \ge 0$ ,  $max_i |\lambda_i(A)| \rightarrow 1$  corresponding to an eigenvector  $\vec{w}$ . From eigenvalue problem,

$$A\vec{w} = max_i |\lambda_i(A)|\vec{w}.$$
(33)

For  $\vec{w}$  such that  $\vec{b}\vec{w} = 1$  with  $\vec{b} = (1,1,\cdots,1)$  we have that,

$$max_i|\lambda_i(A)| = max_i|\lambda_i(A)|\vec{b}\vec{w} = \vec{b}\vec{w} = 1.$$
(34)

In turn, this implies that

$$max_i|\lambda_i(A)| = 1. \tag{35}$$

#### 4. Conclusion

An analytical method based on the decomposition of an admissible perturbation from a set of block diagonal matrices into a block diagonal matrix having identity matrices along the principal diagonal is presented. The main contribution is the computation of the lower bounds of  $\mu$  –values. The proposed methodology is based on the idea of computing spectral and numerical radii of the perturbed matrix. The structured singular value is a tool in control to discuss the stability analysis of the linear system. The lower bounds of structured singular value measure the instability analysis of feedback systems in linear control.

## **Compliance with ethical standards**

## **Conflict of interest**

The authors declare that they have no conflict of interest.

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